Stochastic Thermodynamics and Thermodynamics of Information

Lecture IV: Thermodynamics of Information

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- 1. Reminder
- 2. Information and the second law
- 3. "Information is physical"
- 4. Reversibilities
- 5. Summary
- 6. Reservoirs

- Entropy: Measure of uncertainty of a system?
- 2nd law: Objective? Subjective?
- Manipulation of information (e.g., computing): Does it require dissipation?
- Physical bounds on efficiency of information handling?

Can Stochastic Thermodynamics help?

Entropy: Gibbs and Shannon

• Gibbs formula for the entropy:

$$S^{\rm eq} = -k_{\rm B} \sum_{x} p^{\rm eq} \log p^{\rm eq}$$

• Shannon: entropy of a probability distribution *p*:

$$H(p) = -\sum_{x} p_x \log p_x$$

• Non-equilibrium entropy?

$$S^{\text{n.eq.}} = -k_{\rm B} \sum_x p_x \log p_x$$

- Objection: $S^{\mathrm{n.eq.}}$ is constant for isolated Hamiltonian systems
- We assume it holds for systems in stochastic thermodynamics

- System with energy function $E = (E_x)$, arbitrary probability distribution $p = (p_x)$ in contact with a reservoir at temperature T
- Non-equilibrium free energy

$$\mathcal{F}(p) = \langle E \rangle_p - k_{\rm B} T H(p) = \langle E \rangle_p - T S$$

• Process transforming p from $p^{(0)}$ to $p^{(1)}$:

$$W \ge \mathcal{F}(p^{(1)}) - \mathcal{F}(p^{(0)}) \tag{*}$$

Derivation:

$$0 \leq \Delta S^{\text{tot}} = \Delta S^{(r)} + \Delta S$$
$$= -\frac{Q}{T} + \Delta S = \frac{1}{T} \left(-\Delta \langle E \rangle + W + T \Delta S \right)$$
$$W \geq \Delta \langle E \rangle - T \Delta S$$

Saturating the bound

- Define $\mathcal{H}^{(0)}$: $e^{(F^{(0)} \mathcal{H}^{(0)}_x)/k_B T} = p_x^{(0)}$, $F^{(0)} = -k_B T \log \sum_x e^{-\mathcal{H}^{(0)}_x/k_B T}$
- Perform the sudden transformation (1): $E \longrightarrow \mathcal{H}^{(0)}$. One has Q = 0,

$$W^{(1)} = \Delta \left\langle E \right\rangle_{p^{(0)}} = \sum_{x} p_x^{(0)} \left(\mathcal{H}_x^{(0)} - E_x \right)$$

• Perform a slow (reversible) transformation (2): $\mathcal{H}^{(0)} \longrightarrow \mathcal{H}^{(1)}$, with $\mathcal{H}^{(1)}$: $e^{(F^{(1)} - \mathcal{H}^{(1)}_x)/k_{\rm B}T} = p_x^{(1)}$. One has $W^{(2)} = \Delta \langle \mathcal{H} \rangle - Q$, and, from $\Delta S^{\rm tot} = 0$

$$Q = T \Delta S = T \left(S^{(1)} - S^{(0)} \right)$$

Saturating the bound

- Perform the sudden transformation (3): $\mathcal{H}^{(1)} \longrightarrow E$. One has Q = 0 and

$$W^{(3)} = \Delta \left\langle E \right\rangle_{p^{(1)}} = \sum_{x} p_x^{(1)} \left(E_x - \mathcal{H}^{(1)_x} \right)$$

Therefore

$$W^{(1)} + W^{(2)} + W^{(3)} = \left\langle \mathcal{H}^{(0)} \right\rangle_{p^{(0)}} - \left\langle E \right\rangle_{p^{(0)}} + \left\langle \mathcal{H}^{(1)} \right\rangle_{p^{(1)}} - \left\langle \mathcal{H}^{0} \right\rangle_{p^{(0)}} - T \left(S^{(1)} - S^{(0)} \right) + \left\langle E \right\rangle_{p^{(1)}} - \left\langle \mathcal{H}^{(1)} \right\rangle_{p^{(1)}} = \left\langle E \right\rangle_{p^{(1)}} - TS^{(1)} - \left(\left\langle E \right\rangle_{p^{(0)}} - TS^{(0)} \right)$$

Szilárd's demon revisited





Szilárd's demon revisited

Manipulation:



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Szilárd's demon revisited

Resetting:





The transformation

Initial state: X and Y are independent, $P(X) = P^{eq}(X) = (p_x^{eq})$ Measurement: Introduces correlations between X and Y

$$\begin{split} \Delta \mathcal{F}^{\text{meas}} &= -T \, \Delta S = k_{\text{B}} T \left(I(X:Y) - \Delta H^{\text{meas}}(Y) \right) \\ W^{\text{meas}} &= \langle \mathcal{W} \rangle \geq \Delta \mathcal{F} \geq 0 \end{split}$$

Manipulation: Relax X distribution to $P^{eq}(X)$:

$$W^{\text{extr}} \ge \Delta \mathcal{F}^{\text{man}} \ge -\Delta \mathcal{F}^{\text{meas}} \ge -k_{\text{B}}TI(X:Y)$$

Resetting: Relax *Y* distribution to $P^{(0)}(Y)$:

$$\Delta \mathcal{F}^{\rm res} = -k_{\rm B} T \Delta H^{\rm meas}(Y)$$
$$W^{\rm res} \ge \Delta \mathcal{F}$$

Thus

$$\sum W \geq 0$$

Entropy changes in Y

• If Y is a "clean slate", $H^{(0)}(Y) = 0$, then

$$\Delta S^{\text{meas}} = k_{\text{B}} \left(H^{\text{meas}}(Y) - I(X:Y) \right)$$

which could vanish

In this case

$$W \ge \Delta \mathcal{F}^{\text{res}} = -k_{\text{B}}T \,\Delta H^{\text{meas}}(Y) \ge 0$$

(In Szilard's gas "experiment" $-\Delta H^{\text{meas}}(Y) = \log 2$)

- This is the gist of Landauer's bound: $W \ge k_{\rm B}T\log 2$ for resetting a bit
- $\cdot \ \, {\rm One} \ {\rm can} \ {\rm set} \ P(Y) = P^{\rm eq}(Y) {\rm , \ so \ that} \ \Delta H^{\rm meas}(Y) = 0$

A different initial condition

Initial condition is same as final condition (no resetting):



Information and work balance

• Thus in general, for a memory degree-of-freedom

 $W^{\text{meas}} + W^{\text{res}} \ge k_{\text{B}}TI(X:Y)$

- Inequalities stem from a fluctuation relation for the fluctuating work \mathcal{W} and the fluctuating mutual information $\mathcal{I}_{xy} = -\log(\mathcal{P}(x, y)/(\mathcal{P}(x)p_y))):$ $\left\langle e^{-(\mathcal{W}-\Delta F)/k_{\rm B}T-\mathcal{I}} \right\rangle = 1$ (†)
- More generally, if feedback control is present, manipulation λ depends on measurement outcome y, and one has

$$\left\langle \mathrm{e}^{-(\mathcal{W}-\Delta F)/k_{\mathrm{B}}T} \right\rangle = \gamma$$

where

$$\gamma = \sum_{y} P_{\hat{\lambda}(y)}(y)$$

where $P_{\lambda}(y)$ is the probability of obtaining measurement outcome y with manipulation λ

Proof of (†)

- Assume x is measured at time t_m , $x(t_m) = x_m$ and the outcome is y, then the protocol $\lambda(y)$ is applied
- Crooks: $e^{\Delta S(\boldsymbol{x};y)/k_{\mathrm{B}}} = \mathcal{P}_{\boldsymbol{\lambda}(y)}(\boldsymbol{x})/\mathcal{P}_{\hat{\boldsymbol{\lambda}}(y)}(\hat{\boldsymbol{x}})$, $\forall y$
- $\mathcal{P}(\boldsymbol{x}, y) = p_{y|xm} \mathcal{P}_{\boldsymbol{\lambda}(y)}(\boldsymbol{x})$ $e^{\mathcal{I}} = p_{y|x_m}/p_y$
- Thus

$$\mathrm{e}^{\Delta S(\boldsymbol{x})/k_{\mathrm{B}}+\mathcal{I}_{y|x_{\mathrm{m}}}} = \frac{p_{y|x_{\mathrm{m}}}\mathcal{P}_{\boldsymbol{\lambda}(y)}(\boldsymbol{x})}{\mathcal{P}_{\boldsymbol{\lambda}(y)}(\boldsymbol{\hat{x}})p_{y}}$$

which implies

$$\sum_{y} \int \mathcal{D}\boldsymbol{x} \ \mathcal{P}(\boldsymbol{x}, y) e^{-\Delta S(\boldsymbol{x}, y)/k_{\mathrm{B}} - \mathcal{I}} = \sum_{y} \int \mathcal{D}\boldsymbol{x} \ \mathcal{P}_{\hat{\boldsymbol{\lambda}}(y)}(\hat{\boldsymbol{x}}) p_{y} = 1$$

Exploiting information

Feedback manipulation of a Brownian particle:

TOYABE ET AL., 2010



Exploiting information

Feedback manipulation of a Brownian particle:

Toyabe et al., 2010



Exploiting information

Feedback manipulation of a Brownian particle:

TOYABE ET AL., 2010



- A thermodynamical system handling information:
 - Must possess several "computational" states
 - \cdot These states are long lived
 - Hence it is necessarily non-ergodic
- Computational states correspond to ergodic components of the phase space of the system
- Each is a *macroscopic state*, with its internal energy, entropy, etc.

A toy model



We have to distinguish the degrees-of-freedom:

I: Computational degrees-of-freedom: Here, $X \in \{0, 1\}$ II: Microscopic degrees-of-freedom: Phase-space variables in each ergodic component

III: Heat-bath degrees-of-freedom: At equilibrium at the temperature T

Thermodynamic reversibility: On the whole system (I+II+III)
Logical (computational) reversibility: Connected to entropy change in I
Heat transfer to the bath: Connected to changes in the entropy of I+II

- Thermodynamics implies that the total entropy production $\Delta S^{\rm tot}$ in I+II+III is non-negative
- Transformations are thermodynamically reversible if and only if $\Delta S^{\rm tot}=0$

The different faces of reversibility

- In *Stochastic thermodynamics* we deal with probability distributions *P* (in general non-equilibrium)
- A probability distribution P can be converted into another distribution P' with heat absorption Q if and only if

$$\Delta S + \Delta S^{(\mathbf{r})} = \Delta S^{\text{tot}} \ge 0$$

where ΔS is the change in the system's entropy:

$$\Delta S = k_{\rm B} \left(H(P') - H(P) \right)$$

• As a corollary

$$W \ge \Delta \mathcal{F}$$

• The process $P \longrightarrow P'$ is thermodynamically reversible (i.e., the initial state of the system and the reservoir can be restored) if and only if $\Delta S^{\text{tot}} = 0$

The different faces of reversibility

Computational (or logical) reversibility:

- Let $X \in \{0,1\}^n$ be a collection of bits, and $X \longrightarrow X'$ a computation
- The computation is computationally reversible if X is a one-valued function of $X', \forall X'$
- Examples:
 - NOT (¬) is reversible: $x' = \neg x \Rightarrow x = \neg x'$
 - ERASE (\downarrow) is *not* reversible: $\downarrow 1 = \downarrow 0 = 0$
 - One-bit Boolean functions are not reversible: e.g., AND, OR, XOR...
 - $\cdot\,$ Two-bits mappings such as EXCHANGE can be reversible
- Given a (deterministic) computation $X' = \phi(X)$, ϕ is computationally reversible if and only if the Shannon entropy of any pdf P(X) is equal to $P(\phi(X))$
- Computationally irreversible transformations reduce the entropy of ${\cal P}(X)$

The different faces of reversibility

Reversible erasure

• Heat emission upon erasure of a bit (Landauer bound):

 $-Q \geq k_{\rm B}T\log 2$

• Let us look at Szilard's engine: $W \ge k_{\rm B}T \log 2$:



Any logically irreversible manipulation of information, such as the erasure of a bit or the merging of two computation paths, must be accompanied by a corresponding entropy increase in non-information bearing degrees of freedom of the information processing apparatus or its environment.

Bennett, 2003

But the process can still be thermodynamically reversible if

$$\Delta S^{\text{tot}} = -\frac{Q}{T} + \Delta S^{(\mathsf{I})} + \Delta S^{(\mathsf{II})} = 0$$

JUN ET AL., 2014





p: probability of ending in the right well (p = 1: full erasure)

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• Fluctuating work:

$$\mathcal{W}(\boldsymbol{x}) = \int_0^{ au} \mathrm{d}t \; \dot{\lambda}(t) \partial_{\lambda} U(x(t), \lambda(t))$$
 discretized

• Asymptotic work:

$$\frac{W(\tau)}{k_{\rm B}T} = \frac{W(\infty)}{k_{\rm B}T} + a\tau^{-1}$$

	Asym W	a	χ^2
p = 1	0.71	1.39	8.2
p = 0.5	0.05	1.48	7.5

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- $\cdot\,$ Computation d-o-f contribute as well to the entropy balance
- There is a *subtle* link between computational and thermodynamical reversibility
- There is dissipation in information handling at finite speed: Speed-dissipation tradeoff?

Next: Information handling in biological systems

Mandal-Jarzynski and Barato-Seifert

Bla bla bla

Thank you!

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